Example 8. Consider the Muller automaton over $\Sigma=\{a, b, c\}$ below with $\mathcal{F}=$ $\left\{\left\{q_{c}\right\},\left\{q_{a}, q_{b}\right\},\left\{q_{a}, q_{b}, q_{c}\right\}\right\}$.


This automaton accepts the language
$L=\left\{\alpha \in \Sigma^{\omega} \mid a\right.$ appears infinitely often in $\alpha \Leftrightarrow b$ appears infinitely often in $\left.\alpha\right\}$.

Am example run of the parity automaton constructed according to Definition 17 is shown below.

| $\alpha$ |  | $c$ |  | $a$ |  | $c$ |  | $b$ |  | $b$ |  | $a$ | $a$ | $b$ |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | $\underline{q_{a}}$ | $q_{c}$ | $q_{a}$ | $q_{c}$ | $q_{b}$ | $\underline{q_{b}}$ | $q_{a}$ | $q_{a}$ | $q_{b}$ | $\underline{q_{b}}$ |  |  |  |  |
|  | $q_{b}$ | $q_{a}$ | $\underline{q_{c}}$ | $\underline{q_{a}}$ | $q_{c}$ | $q_{c}$ | $q_{b}$ | $q_{b}$ | $\underline{q_{a}}$ | $q_{a}$ |  |  |  |  |
|  | $q_{c}$ | $\underline{q_{b}}$ | $q_{b}$ | $q_{b}$ | $\underline{q_{a}}$ | $q_{a}$ | $\underline{q_{c}}$ | $q_{c}$ | $q_{c}$ | $q_{c}$ |  |  |  |  |
| hit | 1 | 3 | 2 | 2 | 3 | 1 | 3 | 1 | 2 | 1 |  |  |  |  |
| color | 1 | 6 | 3 | 3 | 6 | 1 | 6 | 1 | 4 | 1 |  |  |  |  |

