

Parametricity and Modular Reasoning

Homework #2

Instructor: Derek Dreyer

Assigned: Tuesday, 6 November 2012
Problem 2: Thursday, 8 November 2012
Due: Tuesday, 13 November 2012

1 Canonical Forms for Sum Types

Prove that if $\vdash e : \tau_1 + \tau_2$, then $(e, e[\tau_1 + \tau_2] \text{inj}_1 \text{inj}_2) \in E[\tau_1 + \tau_2]$.

2 Foldr-build Theorem

The foldr-build theorem underlies short-cut fusion. To state the theorem for System F extended with a base type $\text{list}(\tau)$, we need Church-encoded lists:

$$\begin{aligned}\text{Chlist}(\tau) &\stackrel{\text{def}}{=} \forall \alpha. \alpha \rightarrow (\tau \rightarrow \alpha \rightarrow \alpha) \rightarrow \alpha \\ \text{Cn nil} &\stackrel{\text{def}}{=} \Lambda \alpha. \lambda n. \lambda c. n \\ \text{Chcons} &\stackrel{\text{def}}{=} \lambda x. \lambda L. \Lambda \alpha. \lambda n. \lambda c. c x L.\end{aligned}$$

Prove that if $\vdash f : \text{Chlist}(\tau)$ and $\vdash n : \sigma$ and $\vdash c : \tau \rightarrow \sigma \rightarrow \sigma$, then $\text{fold}(n, c, \text{build } \tau f) \equiv f \sigma n c : \sigma$ where

$$\begin{aligned}\text{build } \tau f &\stackrel{\text{def}}{=} f \text{list}(\tau) \text{nil cons} \\ \text{cons} &\stackrel{\text{def}}{=} \lambda x. \lambda y. \text{cons}(x, y).\end{aligned}$$

(We started a proof in class. See the November 8 scribe notes.)